

# Co-analytic Counterexamples<sup>"</sup> to Marstrand's Projection Theorem<sup>"</sup>

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~~19~~ September 2023

20

Hausdorff dimension



Hausdorff  
dimension

+

Kolmogorov  
complexity



Marstrand  
counterexamples



Kolmogorov complexity

Hausdorff dimension



Hausdorff  
dimension

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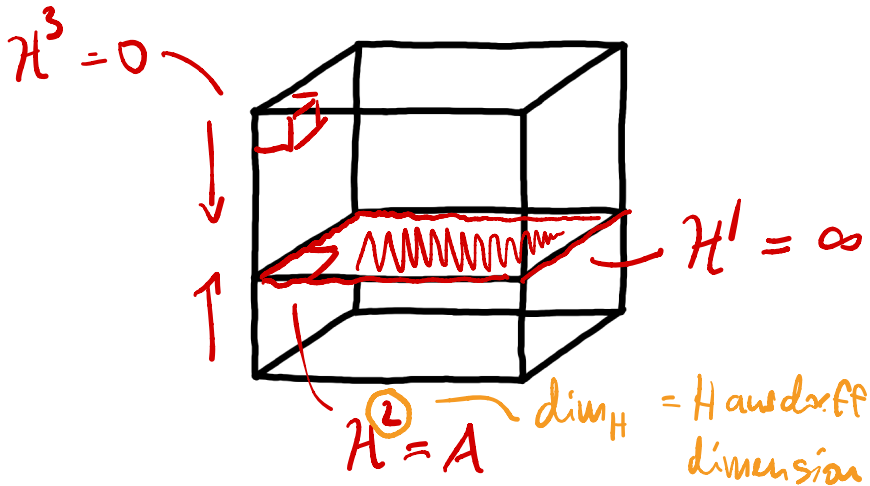


Marstrand  
counterexamples



Kolmogorov complexity

# Hausdorff dimension: motivation



## Definition (Hausdorff dimension)

For  $E \subset \mathbb{R}^n$

$$\dim_H(E) = \sup\{s \mid \mathcal{H}^s(E) = \infty\}$$

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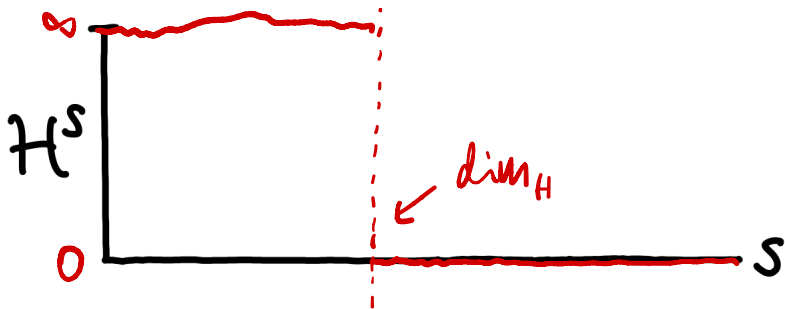
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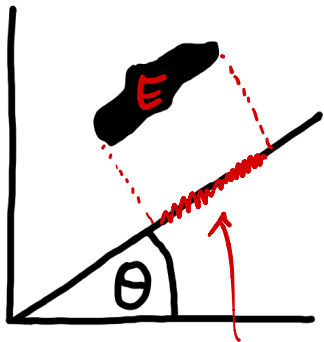


## Lemma

$\dim_H$  is invariant under isometries.

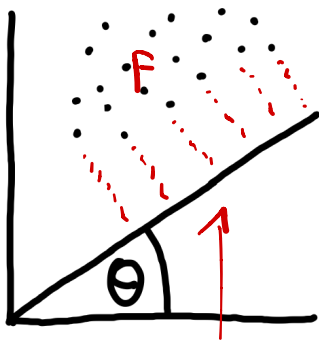
# Marstrand's Projection Theorem

$$\dim_H(E) = \frac{3}{2}$$



$$\dim_H(P(E)) = 1$$

$$\dim_H(F) = \frac{1}{2}$$



$$\dim_H(P(F)) = \frac{1}{2}$$



## Marstrand's Projection Theorem (J. Marstrand (1954))

*Let  $E \subset \mathbb{R}^2$  be analytic. For almost all  $\theta$*

$$\dim_H(p_\theta(E)) = \min\{\dim_H(E), 1\}$$

*where  $p_\theta$  is the orthogonal projection onto the line  $\theta$ .*

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- Also holds if  $\dim_H(E) = \dim_P(E)$  (N. Lutz and Stull 2018)

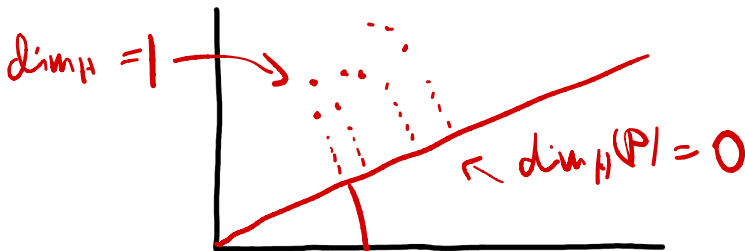
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- Also holds if  $\dim_H(E) = \dim_P(E)$  (N. Lutz and Stull 2018)
- ...but—assuming CH—does *not* hold for all sets (Davies 1979)



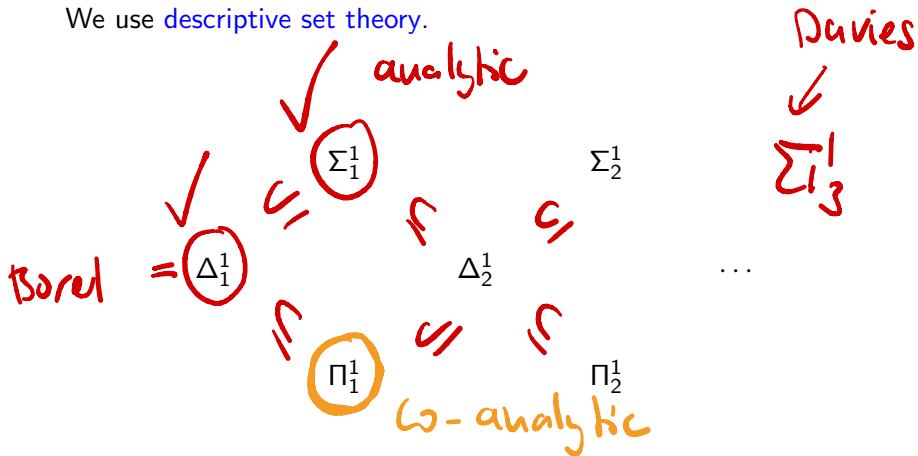
## Question

What is the “simplest” set failing Marstrand's theorem?

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We use descriptive set theory.



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# Strings with long descriptions are complicated

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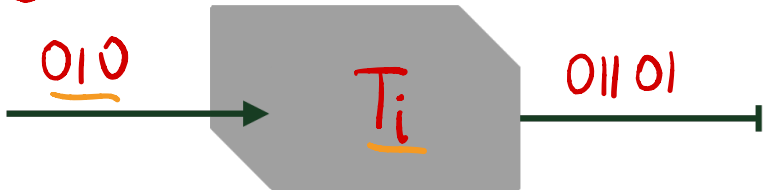
What is the complexity of 01101?

# Strings with long descriptions are complicated

## Question

What is the complexity of 01101?

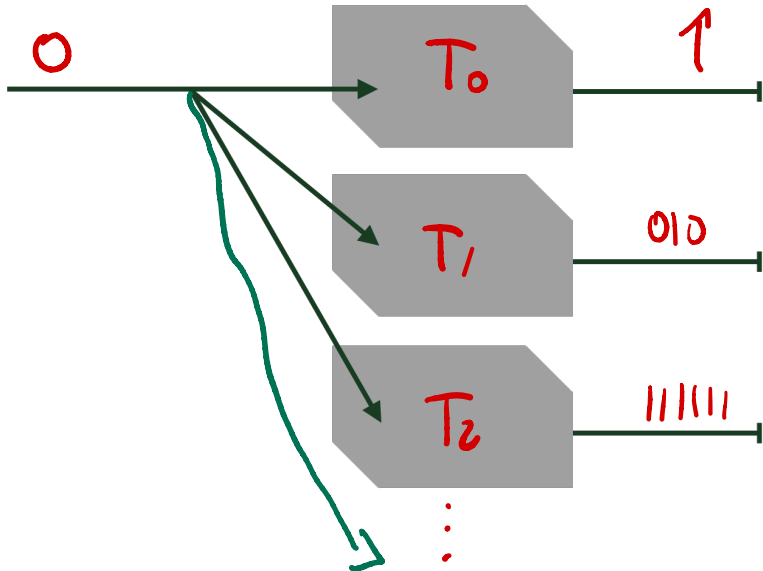
We ask: how long is the **shortest description** of 01101 in a prefix-free machine?



0 x  
1 x  
⋮  
010 ✓

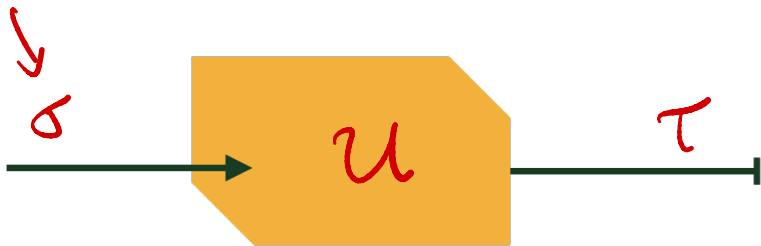
$$C_{T_i}(01101) = \ell(010) = 3$$

# Strings with long descriptions are complicated



Strings with long descriptions are complicated

$l(\sigma) = K(\tau) \leftarrow$  minimal length.



Definition (Levin; Chaitin (1970s))

The **prefix-free complexity** of a string  $\tau$  is  $K(\tau) = C_U(\tau)$ .

$\in 2^{\leq W}$

## Definition (Chaitin; Levin (1970s))

A real  $f \in 2^\omega$  is **Kolmogorov random** if  $K(f[n]) \geq n - c$  for some constant  $c$ .

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The following definitions of randomness are all equivalent!

- 1 **effective open covers**  $\rightarrow$  Martin-Löf (1966)
- 2 **complexity of strings**  $\rightarrow$  Chaitin; Levin (1970s)
- 3 **martingales**  $\rightarrow$  Schnorr (1973)

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So: random reals do not have (long) patterns!



A handwritten binary string in red ink:  $010110101000000000000000000000\dots$ . The first 10 bits are underlined with a yellow bracket. A larger yellow bracket spans from the 11th bit to the end of the string. Below the larger bracket, the text "put 50000 zeros" is written in yellow.

Hausdorff dimension



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## Theorem (J. Lutz; Mayordomo (2003))

There exists  $\dim$  on  $2^\omega$  given by

$$\dim(f) = \liminf_{n \rightarrow \infty} \frac{K(f[n])}{n}.$$

"information density"

*This relativises!*

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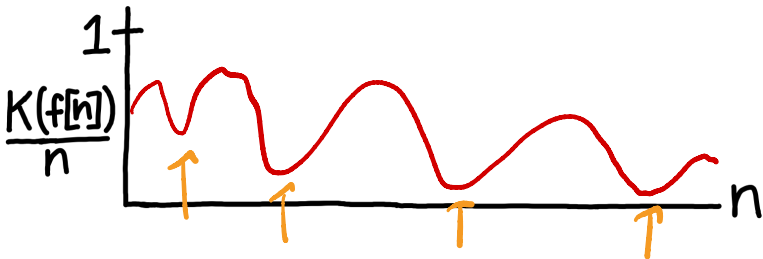
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"information density"  
←

*This relativises!*

- If  $f \in 2^\omega$  is Kolmogorov random then  $\dim(f) = 1$ .
- If  $f \in 2^\omega$  is computable then  $\dim(f) = 0$ .



$x \in \mathbb{R}$  $\in 2^\omega$ 

---

 $x = x_0, x_1, x_2, x_3, x_4, \dots$ 

in binary

global  
property

local  
property

## Point-to-set Principle (J. Lutz, N. Lutz (2018))

For  $E \subset \mathbb{R}^n$  we have

$$\dim_H(E) = \min_{A \in 2^\omega} \sup_{x \in E} \dim^A(x).$$

- Extension of result due to Hitchcock (2003) for lightface  $\Pi_1^0$  classes on  $2^\omega$ .

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**Marstrand  
counterexamples**

Kolmogorov complexity



## Marstrand's theorem (special case)

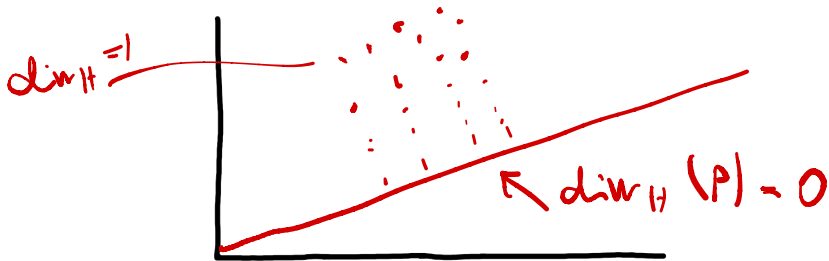
For every analytic  $E \subset \mathbb{R}^2$  for which  $\dim_H(E) = 1$  we have  $\dim_H(p_\theta(E)) = 1$  for almost all  $\theta$ .

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## Theorem (R.)

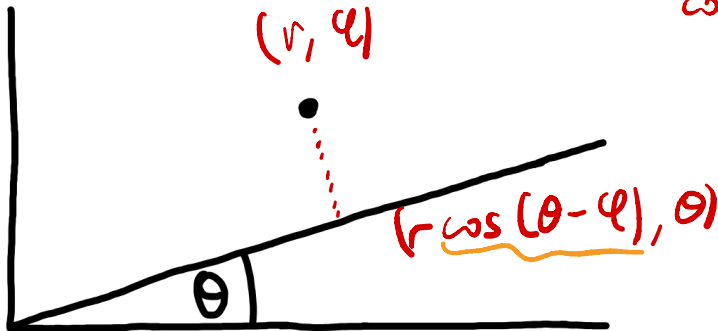
( $V=L$ ) There exists a **co-analytic**  $E \subset \mathbb{R}^2$  such that  $\dim_H(E) = 1$  and  $\dim_H(p_\theta(E)) = 0$  for **all**  $\theta$ .



point-to-set principle

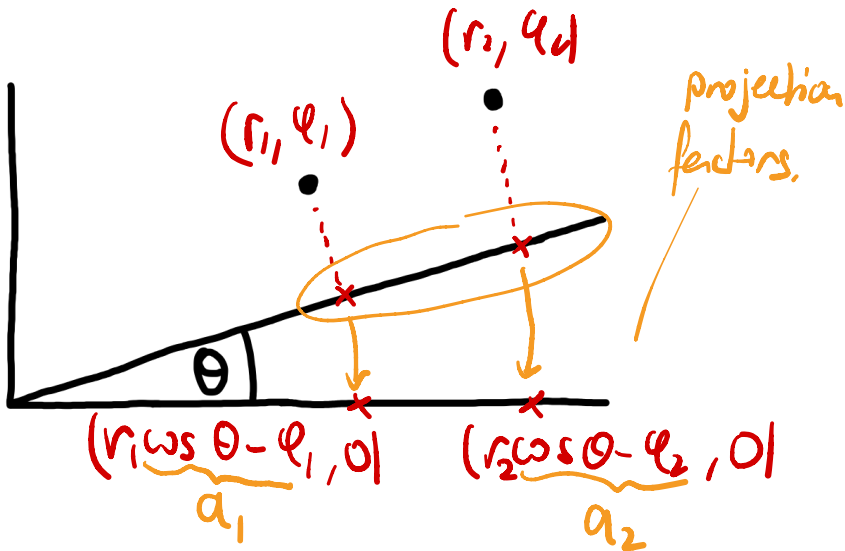
Recall:  $\dim_H(E) = \min_{A \in 2^\omega} \sup_{x \in E} \dim^A(x)$

in polar coordinates





Recall:  $\dim_H$  is invariant under isometries.



## How do we **control dimension**?

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## Proof.

Let  $A \in 2^\omega$ . Take  $\theta$  random relative to  $A$ .

so  $\dim^A(\theta) = 1$

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$$\dim^A(\underline{r}, \theta) \geq \dim^A(\theta) = 1.$$

Since random wrt. to  $A$ .

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$A$  is arbitrary, so PTS completes the argument. □

# How do we **construct co-analytic sets**?

By recursion!

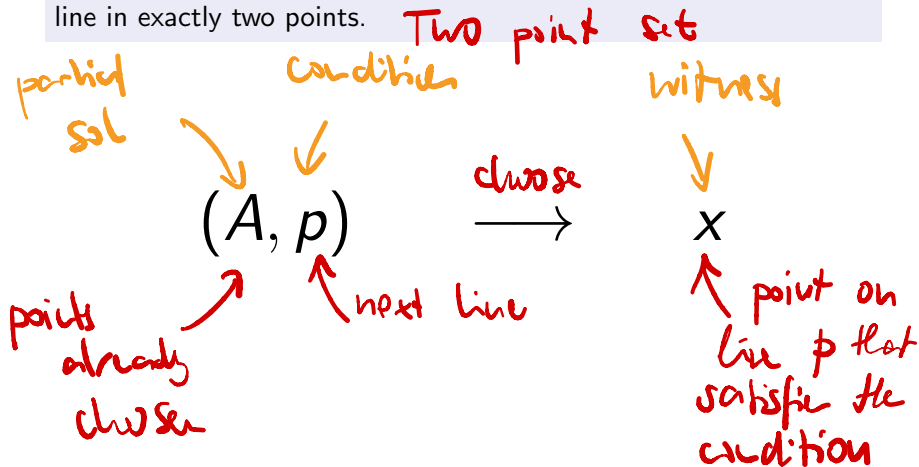


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By recursion! How we normally do recursion:

Example (Mazurkiewicz 1914)

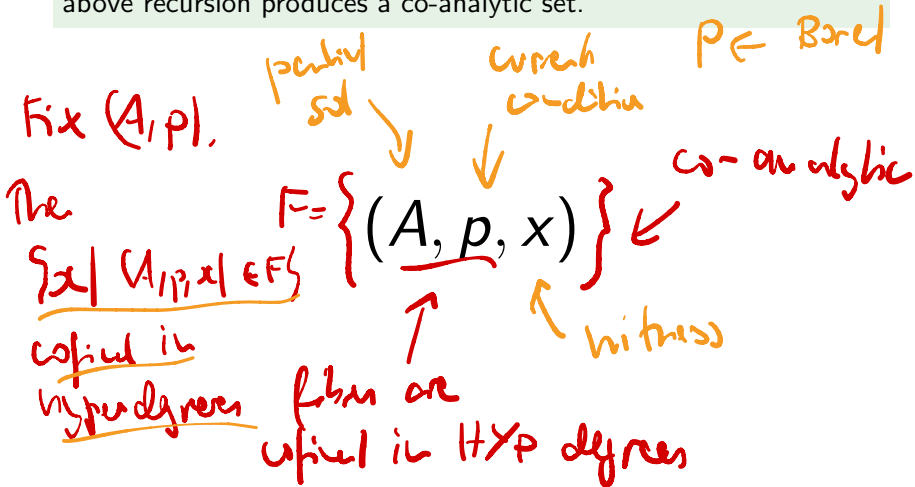
(AC) There exists a subset  $B \subset \mathbb{R}^2$  that intersects every straight line in exactly two points.



# How do we construct co-analytic sets?

Theorem (Z. Vidnyánszky (2014))

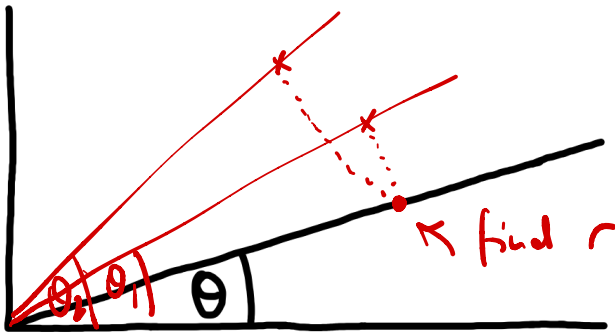
$(V=L)$  Under certain conditions on the set of triples  $(A, p, x)$ , the above recursion produces a co-analytic set.



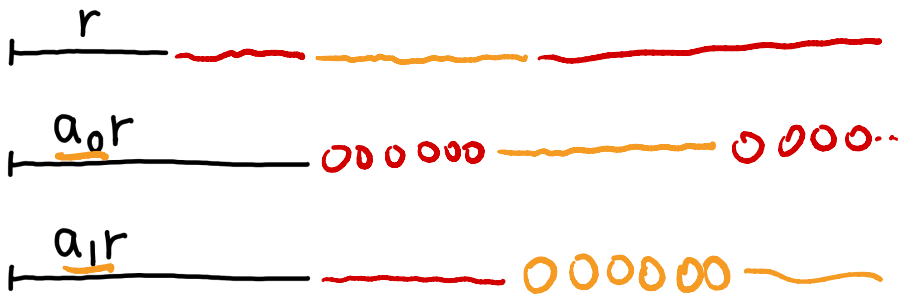
## Constructing $E$ by recursion

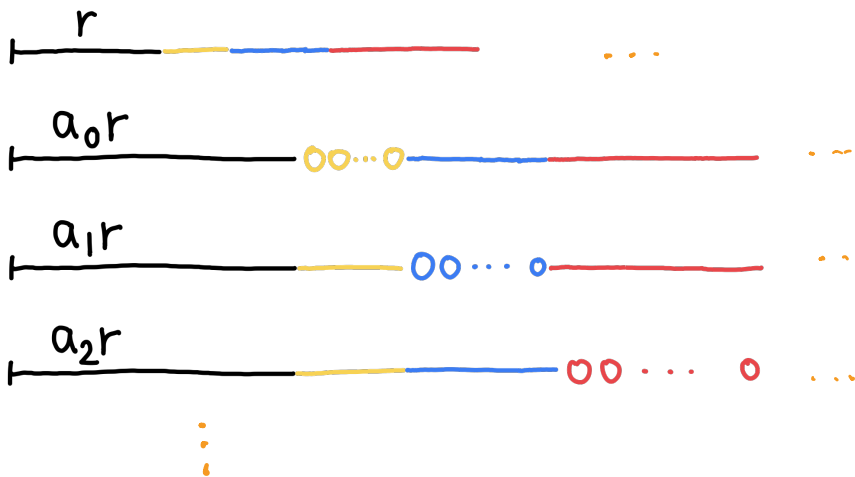
- do recursion on all lines  $\theta$  through the origin
- at step  $\theta$ , take all previous lines  $\theta_0, \theta_1, \theta_2, \dots$
- find  $r$  so that  $\dim(p_{\theta_i}(r, \theta)) = \dim(a_i r) = 0$
- enumerate  $(r, \theta)$  into  $E$

We assumed  
 $V = L + CH$

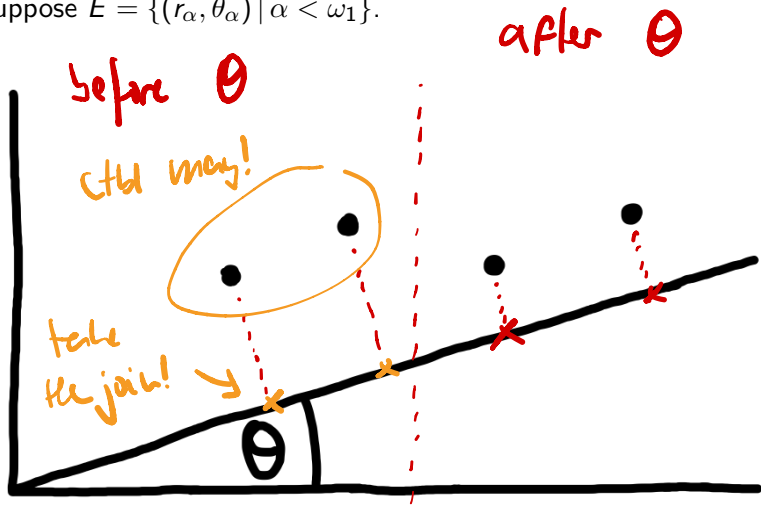


# Constructing $r$





Suppose  $E = \{(r_\alpha, \theta_\alpha) \mid \alpha < \omega_1\}$ .



## Theorem (R.)

( $V=L$ ) For every  $\epsilon \in [0, 1)$  there exists a *co-analytic*  $E_\epsilon \subset \mathbb{R}^2$  such that  $\dim_H(E_\epsilon) = 1 + \epsilon$  and  $\dim_H(p_\theta(E_\epsilon)) = \epsilon$  for all  $\theta$ .

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PTS for packing dimension (J. Lutz, N. Lutz (2018))

$$\dim_P(E) = \min_{A \in 2^E} \sup_{x \in E} \text{Dim}^A(x)$$

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- Is there a co-analytic set failing Marstrand's theorem that is not thin?

Thank you