

Genetic Programming Hyper-heuristic for Stochastic Team Orienteering Problem with Time Windows

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Abstract—This paper investigates the stochastic team orienteering problem with time windows, which is a well known problem to model personalised tourist trip design. Specifically, we consider the stochastic visit duration, which may make preplanned trip infeasible. Existing studies focus on optimising robust solutions in advance, which is not effective in adjusting the subsequent trip in real time. Decision making policies, on the other hand, are effective heuristics to this end. However, it is very challenging to manually design effective policies. In this paper, we investigate automatically evolving policies for the stochastic team orienteering problem with time windows by genetic programming hyper-heuristics. We designed novel problem-specific features for the terminal set, and a meta-algorithm for fitness evaluation. Furthermore, we developed two look-ahead features that can provide more fruitful information than the basic features for real-time decision making. The experimental studies showed that the proposed genetic programming hyper-heuristic can evolve policies that are much better than the manually designed policies. In addition, it seems that the look-ahead features are not so effective when directly included in the terminals. This suggests the requirement of more intelligent ways of incorporating look-ahead information.

I. INTRODUCTION

Personalised Tourist Trip Design Problem (TTDP) [1] is an important and challenging problem to solve. It is to design a (multi-day) trip, which is a sequence of visits to the Places-of-Interests (POIs), for a tourist under a given budget, so that the tourist’s satisfaction of the trip is maximised. A typical objective is to maximise the total score (reflecting the satisfaction of the tourist) of the POIs visited during the trip.

TTDP can be formulated as an orienteering problem [2], which contains a variety of models. In this paper, we focus on the Team Orienteering Problem with Time Windows (TOPTW) [3] due to its closeness to the reality. TOPTW aims to design a multi-day (“teamed”) trip where the visit of each POI must start within its time window. First, a tourist usually stays in a city/region for multiple days, and thus it is more reasonable to design a multi-day trip instead of a single-day trip. Second, many POIs (e.g. museums, parks, special events) have opening hours, and it is impossible to visit them outside their opening hours. TOPTW exactly considers these two real-world factors.

In this paper, we consider another important factor in addition to the above two factors, i.e. the stochastic visit durations. In real world, it is common that the actual visit

durations are different from the expected values. The tourist may find a POI very attractive, and stay there longer than expected, while dislike another POI and leave much earlier than expected. In this case, the planned trip may become infeasible, as a longer stay of a POI may result in arriving some subsequent POI later than its time window. Therefore, in the *stochastic* TOPTW, it is necessary to adjust/plan the subsequent trip in real time.

There have been a number of studies for TOPTW (e.g. [3], [4], [5], [6], [7], [8]). However, none of them considered the stochastic visit duration, and the preplanned trip designed by existing studies may become infeasible in real time. As a result, the existing studies are not directly applicable to solving stochastic TOPTW.

Reactive decision making policies such as dispatching rules for job shop scheduling have shown to be very effective to make real-time response in a dynamic/stochastic environment. However, manually designing effective such policies is time-consuming and demands domain expertise. To address this issue, Genetic Programming Hyper-Heuristic (GPHH) has been successfully applied to automatically design (evolve) effective decision making policies for dynamic scheduling [9], [10] and routing [11], [12].

To the best of our knowledge, our work is the first time to consider evolving decision making policies for stochastic orienteering problems. This paper aims to propose a GPHH algorithm to evolve effective policies that make real-time decisions for stochastic TOPTW. More specifically, it contains the following research objectives.

- To design the *meta-algorithm*, i.e. the algorithm to generate a feasible solution for stochastic TOPTW given an instance and a policy.
- Design the features to be used by the GPHH to evolve effective policies.
- Analyse the structure of the evolved policies and understand the importance of the designed features in effective policies.

The rest of the paper is organised as follows. Section II gives the background introduction, including the problem description and related work. Then, the proposed GPHH is described in III. Then, the experimental studies and analysis

are carried out in Section IV. Finally, the conclusions and future work are given in Section V.

II. BACKGROUND

A. Problem Description

In TOPTW, a set of n POIs $\mathbf{P} = \{1, \dots, n\}$ are given. Each POI $p \in \mathbf{P}$ has a score $s(p)$, a visit duration $d(p)$ and a time window $[o(p), c(p)]$. A starting point $p_s \in \mathbf{P}$ and an ending point $p_e \in \mathbf{P}$ is also given. For each pair of POIs (p_i, p_j) , the travel time $t(p_i, p_j)$ to travel from p_i to p_j is known. The time budget is given by m days, where each day has a starting time T_s and an ending time T_e . Then, the problem aims to design a set of sequences of POIs so that the total score of the visited POIs is maximised, and the following constraints are satisfied.

- 1) There are m sequences of POIs, each for a day.
- 2) Each sequence departs from p_s at time T_s , and returns to p_e no later than T_e .
- 3) Each POI is visited at most once.
- 4) The visit of p_i starts no earlier than o_i and no later than c_i . That is, if the tourist arrives p_i earlier than o_i , s/he has to wait until o_i to start the visit.

Assuming that a solution X is represented by $X = \{X_1, \dots, X_m\}$, where $X_k = (x_{k1}, x_{k2}, \dots, x_{kL_k})$, then the problem can be modelled as follows.

$$\max f(X) = \sum_{k=1}^m \sum_{l=1}^{L_k} s(x_{kl}), \quad (1)$$

$$s.t. \quad x_{k1} = p_s, \quad x_{kL_k} = p_e, \quad k = 1, \dots, m \quad (2)$$

$$v(x_{k1}) = T_s, \quad v(x_{kL_k}) \leq T_e, \quad k = 1, \dots, m \quad (3)$$

$$v(x_{k(l+1)}) \geq v(x_{kl}) + d(x_{kl}) + t(x_{kl}, x_{k(l+1)}), \quad (4)$$

$$o(x_{kl}) \leq v(x_{kl}) \leq c(x_{kl}), \quad l = 2, \dots, L_k - 1, \quad (5)$$

$$x_{k_1 l_1} \neq x_{k_2 l_2}, \quad k_1 \neq k_2 \text{ or } l_1 \neq l_2, \quad (6)$$

$$x_{kl} \in \mathbf{P}, \quad (7)$$

where $v(x_{kl})$ stands for the visit starting time of x_{kl} . Eq. (1) aims to maximise the total score of the visited POIs. Eq. (2) means that all the daily trips depart at p_s and return to p_m . Eq. (3) specifies that each daily trip starts at T_s and finishes no later than T_e . Eqs. (4) and (5) indicate that the visit of each POI cannot start before its arrival, and has to be within its time window. Eq. (6) means that each POI is visited no more than once. Finally, Eq. (7) is the domain constraint.

In Stochastic TOPTW, the visit duration $d(p)$ is a random variable instead of a deterministic value. In this case, the objective becomes maximising the *expected* total score.

B. Related Work

Orienteering Problem has been studied for many years [13], [2]. So far, researchers have considered many different variants such as multi-objective orienteering [14], time-dependent orienteering [15], team orienteering [16], orienteering with time windows [3] and mixtures of them [6], and have proposed effective solution methods such as mathematical programming [17], individual-based search (e.g. iterated local

search [3], simulated annealing [7] and tabu search [18]) and population-based search (e.g. genetic algorithms [15], [8] and ant colony optimisation [19], [20]). However, these existing studies mainly focus on deterministic problems, where the information is fully known in advance, and will not change during the execution of the solution (i.e. during the process of visiting the POIs based on the planned trip).

There are a few works investigating stochastic orienteering problems. In [21], [22], a sampling-based method was proposed to approximate the expected objective value under stochastic travel and service times. A bi-objective genetic algorithm was proposed for orienteering problem with stochastic profit [23]. Efficient approximation algorithms were proposed for orienteering problem with stochastic service times [24], [25]. A variable neighbour search [26] was proposed for orienteering problem with stochastic travel and service times. However, all the existing studies focus on finding a robust preplanned *solution* that is expected to work well in all possible realised situations. Our work is the first attempt to find a *policy* that effectively reacts to real-time environment changes.

GPHH for automated design of policies has been successfully applied to a number of challenging combinatorial optimisation problems such as scheduling and vehicle routing, and many works have been done to investigate representation [27], [28], feature selection [29], [30] and generalisation [31]. These studies demonstrate the effectiveness of GPHH to evolve policies for dynamic/uncertain combinatorial optimisation problems. Therefore, in this paper, we will investigate the possibility of evolving policies for stochastic TOPTW by GPHH.

III. GENETIC PROGRAMMING HYPER-HEURISTIC FOR STOCHASTIC TEAM ORIENTEERING PROBLEM WITH TIME WINDOWS

This section describes the GPHH algorithm proposed for evolving policies for stochastic TOPTW. The framework is given in Algorithm 1, which is a standard evolutionary algorithm framework. However, the individual representation (especially the terminal set of GP) and fitness evaluation are problem-specific, and will be described next.

Algorithm 1: The framework of GPHH

- 1 Randomly initialise a population of policies;
 - 2 **while** *Stopping criteria not met* **do**
 - 3 *Evaluate* the policies in the current population;
 - 4 Generate an offspring population by applying crossover/mutation/reproduction operators;
 - 5 Select the policies from the current and offspring populations to form the new population;
 - 6 **end**
-

A. Individual Representation

The proposed GPHH uses the traditional tree-based representation, and an individual (policy) is represented as a tree. It is essentially a priority function of the features of a POI.

Given any state during the decision making process, the policy calculates the priority of each candidate POI, and selects the one with the highest priority value to visit next. Figure. 1 shows an example of the policy “Score/TravelTime”. In other words, the policy tends to select the POI with the highest score and shortest travel time from the current place to visit next.

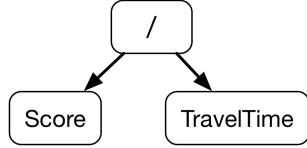


Fig. 1. An example of the policy Score/TravelTime.

The terminal and function sets are critical to the performance of GP [9]. In the proposed GP, the function set is set to $\{+, -, \times, /, \min, \max\}$, where all the operators take two arguments. Note that the “/” operator is protected, i.e. it returns 1 if divided by zero.

The selection of the terminals is problem dependent. It has been shown that time-invariant features are more effective in evolving dispatching rules for dynamic job shop scheduling [32]. Therefore, we design time-invariant features, i.e. the features relative to the current time, for stochastic TOPTW as well. More specifically, the features included in the terminal set of GPHH are described in Table I. For calculating the features involving the random duration $d(p)$, the expected duration is used.

TABLE I
THE FEATURES TO BE USED IN THE TERMINAL SET OF GPHH.

| Notation | Description | Calculation |
|----------|---------------------------------|---|
| SCORE | score of the POI | $s(p)$ |
| DUR | duration of the visit | $E(d(p))$ |
| TO | time to the opening time | $o(p) - \text{currTime}$ |
| TC | time to the closing time | $c(p) - \text{currTime}$ |
| TA | time to arrive the POI | $t(\text{currPlace}, p)$ |
| TR | time to return to the end point | $t(p, p_e)$ |
| TSV | time to start the visit | $\max\{\text{TO}, \text{TA}\}$ |
| TFV | time to finish the visit | $\text{TSV} + E(d(p))$ |
| SL | slack | $\text{TC} - \text{TA}$ |
| RemT | remaining time budget | $\text{remainDays} \cdot (T_e - T_s) + (T_e - \text{currTime})$ |

From the table, one can see that some features seem to be redundant as they can be represented by the combination of other features (e.g. $\text{TFV} = \text{TSV} + \text{DUR}$). However, previous studies [9] have shown that directly keeping promising (constructed) features may be helpful to GP evolution. Therefore, we keep all the features in Table I in the terminal set.

In addition to the above features, two look-ahead features are designed as well. They are described as follows.

- 1) MNS: the maximal score among the *feasible* next POIs.
- 2) ANS: the average score of the *feasible* next POIs.

For a candidate POI p , a feasible next POI p' must satisfy the following constraint:

- p' is unvisited and $p' \neq p$.

- After visiting p , p' can be arrived before its closing time, i.e. $\text{currTime} + \text{TFV}(p) + t(p, p') \leq c(p')$.
- After visiting p and p' , the trip can return to the end point by T_e , i.e. $\max\{\text{currTime} + \text{TFV}(p) + t(p, p'), o(p')\} + d(p') + t(p', p_e) \leq T_e$.

Obviously, the calculations of MNS and ANS are much more time consuming than the features in Table I, as they require enumeration of all the unvisited POIs. However, they provide look-ahead information which can form more effective policies. In the experimental studies, we will investigate the effectiveness of the look-ahead features by comparing the GPHH with and without MNS and ANS.

B. Fitness Evaluation and Meta-Algorithm

The fitness evaluation of a GP individual is different from the traditional evaluation of a solution, since a GP individual is a policy rather than a solution. To evaluate a policy, one needs a set of sampled TOPTW instances (i.e. *training instances* $\mathcal{I}_{\text{train}}$) and a *meta-algorithm* ($\text{solution}()$) that generates a feasible solution given any TOPTW instance and a policy. Then, one can approximate the fitness of a policy as the average total score of the solutions generated for training instances. That is, for a policy pol ,

$$\text{fit}(pol) = \frac{1}{|\mathcal{I}_{\text{train}}|} \sum_{I \in \mathcal{I}_{\text{train}}} \text{score}(\text{solution}(I, pol)). \quad (8)$$

The meta-algorithm is described in Algorithm 2. It is a event-driven decision process, where the events are stored in a priority queue. An event is triggered whenever a visit of a POI is finished, and the policy is applied to select the next POI to visit. At first, the decision process state (i.e. current day and time, current place and tour, and unvisited POIs) and the event queue are initialised (lines 1–3). Then, whenever a POI is visited, all the feasible unvisited POIs are examined. In line 6, an unvisited POI p is feasible if it satisfies the following constraints.

- p can be arrived before its closing time, i.e. $t + t(p^*, p) \leq c(p)$.
- After visiting p , the trip can return to the end point by T_e , i.e. $\max\{t + t(p^*, p), o(p)\} + d(p) + t(p, p_e) \leq T_e$.

Note that in line 22, the duration $d_{\text{sample}}(p_{\text{next}})$ is the actual (sampled) value. However, when examining the feasibility (line 6) and calculating the priority (line 14) of each unvisited POI, the actual visit duration is unknown since the POI is not visited yet. Therefore, the expected value $E(d(p))$ is used in these two situations.

C. Overall Algorithm

In this subsection, we give a more detailed description to the GPHH for stochastic TOPTW. The pseudo code is given in Algorithm 3. At each generation, the offspring population is generated by three genetic operators, i.e. crossover, mutation and reproduction. In addition, some elites are inherited from the previous population to improve convergence.

Note that in line 3, only one training instance is used ($|\mathcal{I}_{\text{train}}| = 1$). However, in each generation, a different sampled

Algorithm 2: solution(\cdot, \cdot)

Input: A TOPTW instance I , a policy pol .
Output: A feasible solution $X = \{X_1, X_2, \dots, X_m\}$.

```

// initialise decision process
1  $day \leftarrow 0, t \leftarrow T_s, p_c \leftarrow p_s, X_{day} \leftarrow (p_s);$ 
2  $\Omega \leftarrow \mathbf{P} \setminus \{p_s, p_e\};$  // unvisited POIs
3  $queue \leftarrow \{(p_s, t)\};$  // initial event queue
4 while  $queue \neq \emptyset$  do
    // trigger the next event
5      $(p^*, t^*) \leftarrow \text{poll}(queue);$ 
6     Update the feasible unvisited POIs  $\Omega' \subseteq \Omega;$ 
7     if  $\Omega' = \emptyset$  then
8          $X_{day} \leftarrow (X_{day}, p_e);$  // return to  $p_e$ 
9          $day \leftarrow day + 1;$  // go to the next day
10        if  $day = m$  then Return  $X;$ 
11         $X_{day} \leftarrow (p_s);$  // open a new tour
12         $queue \leftarrow queue \cup (p_e, T_s);$ 
13    end
14    Calculate the priority value  $pol(p)$  of each  $p \in \Omega';$ 
15     $p_{next} \leftarrow \arg \max_{p \in \Omega'} pol(p);$ 
16     $\Omega \leftarrow \Omega \setminus p_{next};$ 
17     $t_{arr} \leftarrow t + t(p^*, p_{next});$  // arrival time
18     $t_{start} \leftarrow \max\{t_{arr}, o(p_{next})\};$  // visit start time
19    if  $t_{start} > c(p_{next})$  then
    // violate the time window, simply
    // skip the visit
20     $queue \leftarrow queue \cup (p_{next}, t_{start});$ 
21    else
22    sample the actual duration  $d_{\text{sample}}(p_{next});$ 
23     $t \leftarrow t_{start} + d_{\text{sample}}(p_{next});$ 
24     $X_{day} \leftarrow (X_{day}, p_{next});$ 
25     $queue \leftarrow queue \cup (p_{next}, t);$ 
26    end
27 end

```

instance is generated (i.e. using a different random seed) by resampling the random duration variables. Such “instance rotation” has been shown to lead to much better generalisation of GPHH in other problems (e.g. [33], [11]).

IV. EXPERIMENTAL STUDIES

A. Experiment Settings

The commonly used c^*_100 , r^*_100 and rc^*_100 [34], c^*_200 , r^*_200 and rc^*_200 [20] and pr^* [34], [20] are used in our experimental studies. These test instances were extended from vehicle routing instances [35], [36], in which the number of POIs ranges from about 50 to almost 300. Due to page limit, we selected 5 instances for each dataset (e.g. c101~c105) in the experiments.

For each TOPTW instance, the stochastic version is generated by replacing each deterministic $d(p)$ value with a random variable $D(p) \sim \mathcal{N}(d(p), \sigma \cdot d(p))$, where σ is called the *uncertainty level* of the random variable. Here, the uncertainty level is set to $\sigma = 0.2$ as a rule of thumb. That is, the standard deviation of each duration is 20% of its mean value. For each instance, the versions with $m = 1$ (single-day trip) and $m = 3$ (3-day trip) are tested.

Table II shows the parameter setting of the GPHH proposed in this paper. The algorithm was implemented in Java using

Algorithm 3: Pseudo code of GPHH for stochastic TOPTW

```

1 Initialise a population  $pop$  of policies by ramp-half-and-half;
2 while stopping criteria is not met do
3     Sample an instance by sampling a value for all the
    random variables;
4     Evaluate individuals in  $pop;$ 
5      $pop' \leftarrow \text{elites}(pop);$ 
6     while  $|pop'| < |pop|$  do
7         Randomly sample  $r \in U(0, 1);$ 
8         if  $r < P_x$  then
    // Do crossover
9         Select two parents  $pol_1$  and  $pol_2$  from  $pop$  by
    tournament selection;
10        Apply tree-based crossover to  $pol_1$  and  $pol_2$  to
    generate two offsprings  $pol'_1$  and  $pol'_2;$ 
11         $pop' \leftarrow pop' \cup \{pol'_1, pol'_2\};$ 
12        else if  $r < P_x + P_m$  then
    // Do mutation
13        Select  $pol \in pop$  by tournament selection;
14        Apply tree-based mutation to  $pol$  to generate an
    offspring  $pol';$ 
15         $pop' \leftarrow pop' \cup pol';$ 
16        else
    // Do reproduction
17        Select  $pol \in pop$  by tournament selection;
18         $pop' \leftarrow pop' \cup pol;$ 
19        end
20    end
21     $pop \leftarrow pop';$ 
22 end
23 Return best individual in  $pop;$ 

```

the ECJ library [37] and run on Linux OS with Intel(R) Core(TM) i7 CPU @3.60GHz. The GPHH with and without look-ahead features (i.e. MNS and ANS described in Section III-A) are compared. The two GPHH versions are named BasicGP and LA-GP respectively. Each compared algorithm was run 30 times independently, and the results were compared statistically. In addition, two manually designed policies are included in the comparison as a reference. One is the Score-Over-Slack (SOS, with priority function SCORE/SL) policy, and the other is the Score-Over-Time (SOT, with priority function SCORE/TA) policy.

TABLE II
THE PARAMETER SETTING OF THE GPHH.

| Parameter | Value |
|---------------------------|---------------------|
| Population size (popsize) | 1024 |
| Elitism | 10 best individuals |
| Tournament selection size | 7 |
| Maximal depth | 8 |
| Crossover rate | 80% |
| Mutation rate | 15% |
| Reproduction rate | 5% |
| Number of generations | 51 |

For each stochastic TOPTW instance, all the compared policies will be evaluated on a test set with 500 randomly sampled instances. That is, the test fitness of a policy is

calculated as

$$\text{fit}_{\text{test}}(\text{pol}) = \frac{1}{|\mathcal{I}_{\text{test}}|} \sum_{I \in \mathcal{I}_{\text{test}}} \text{score}(\text{solution}(I, \text{pol})), \quad (9)$$

where $|\mathcal{I}_{\text{test}}| = 500$. With such a large sample size, one can achieve a precise approximation of the expected performance over all the possible samples.

B. Experimental Results

Table III shows the test fitness values of BasicGP, LA-GP and the two manually designed policies (SOS and SOT) on the test instances with $m = 1$. In addition, BasicGP is compared with LA-GP statistically using Wilcoxon’s rank sum test with significance level $\alpha = 0.05$. If BasicGP is significantly better than LA-GP, then the corresponding result is marked with “(+)”. If BasicGP is significantly worse than LA-GP, then it is marked with “(-)”.

From Table III, it is obvious that the policies evolved by BasicGP and LA-GP are both much better than the manually designed policies. On some instances (e.g. rc101, rc102, rc105 and r203), the average total score of the GP-evolved policies are even more than twice as the total score obtained by the manually designed policies. This demonstrates that the stochastic TOPTW is so challenging that it is hardly possible to manually design effective policies. In addition, the performance of SOS and SOT can be much different on the same instance (e.g. rc105 and c204). Note that the slack includes more information than the time to arrive (i.e. $SL = TC - TA$). The advantage of SOT over SOS on a decent number of instances (e.g. rc101, rc105, pr03 and pr13) indicates that using more advanced features like slack in manually designed policies may not always be helpful.

When comparing between BasicGP and LA-GP, one can be found that they perform statistically similar in almost all the instances with $m = 1$ (except pr01, where BasicGP significantly outperformed LA-GP). This indicates that simply including the look-ahead features into the terminal set of GP does not seem to improve the test performance of the GP-evolved policies.

Table IV shows the corresponding results for the instances with $m = 3$, i.e. 3-day trips. One can observe similar patterns as those in $m = 1$. The GP-evolved policies are much better than the manually designed policies. In addition, BasicGP performed statistically the same as LA-GP on most instances (LA-GP significantly outperformed BasicGP on 3 instances with “(-)”, while defeated by BasicGP on 4 instances with “(+)”.)

C. Further Analysis

First, we investigate the generalisation of the proposed GPHH algorithm. To this end, we arbitrarily selected eight instances (c103, r101, rc101, rc105, c203, r202, pr03, and pr14 with $m = 1$) and plot the convergence curves of the test performance of BasicGP and LA-GP on these instances in Fig. 2. Other instances show similar patterns. The ribbons are the standard deviation. From the figure, one can see that

TABLE III

THE TEST FITNESS OF THE COMPARED GPHH METHODS ON THE INSTANCES WITH $m = 1$. IF BASICGP IS STATISTICALLY SIGNIFICANTLY BETTER THAN LA-GP UNDER WILCOXON’S RANK SUM TEST WITH $\alpha = 0.05$, THEN THE CORRESPONDING RESULT IS MARKED WITH “(+)”. IF BASICGP IS STATISTICALLY SIGNIFICANTLY WORSE THAN LA-GP, THEN THE CORRESPONDING RESULT IS MARKED WITH “(-)”.

| Instance | SOS | SOT | BasicGP | LA-GP |
|----------|--------|--------|-----------------|----------------|
| c101 | 187.16 | 159.74 | 292.44(10.25) | 289.33(13.79) |
| c102 | 205.96 | 183.52 | 338.90(5.68) | 335.62(7.57) |
| c103 | 264.46 | 245.84 | 375.89(8.26) | 377.39(4.26) |
| c104 | 257.64 | 257.40 | 387.61(6.43) | 386.82(7.75) |
| c105 | 192.32 | 171.42 | 323.02(8.56) | 324.66(7.18) |
| r101 | 98.24 | 74.38 | 189.52(4.25) | 187.59(11.06) |
| r102 | 121.10 | 175.03 | 273.66(6.82) | 275.10(3.10) |
| r103 | 121.10 | 93.10 | 281.19(6.18) | 280.32(8.32) |
| r104 | 149.09 | 128.04 | 288.08(11.90) | 285.68(11.88) |
| r105 | 100.97 | 88.00 | 224.94(6.02) | 221.98(10.18) |
| rc101 | 68.07 | 98.00 | 203.83(9.05) | 202.91(9.81) |
| rc102 | 66.41 | 102.29 | 233.54(7.20) | 234.81(9.03) |
| rc103 | 91.52 | 100.56 | 244.39(5.31) | 241.93(10.36) |
| rc104 | 129.05 | 129.68 | 253.77(8.32) | 254.80(7.03) |
| rc105 | 57.47 | 104.37 | 217.59(16.48) | 220.19(14.61) |
| c201 | 210.00 | 182.16 | 822.96(10.20) | 824.42(13.19) |
| c202 | 268.52 | 183.50 | 867.84(12.59) | 867.96(14.76) |
| c203 | 435.94 | 233.66 | 916.61(11.48) | 915.46(10.31) |
| c204 | 469.00 | 197.60 | 933.20(9.23) | 929.83(6.67) |
| c205 | 273.72 | 188.40 | 861.22(9.63) | 861.09(8.38) |
| r201 | 211.00 | 144.69 | 720.94(18.50) | 720.80(14.35) |
| r202 | 309.73 | 257.08 | 840.37(15.17) | 836.37(12.62) |
| r203 | 325.51 | 257.08 | 923.57(9.65) | 924.27(12.34) |
| r204 | 522.80 | 307.24 | 1004.91(16.66) | 1006.66(10.45) |
| r205 | 226.92 | 240.40 | 831.81(12.11) | 828.60(15.25) |
| rc201 | 182.99 | 117.49 | 729.19(20.24) | 734.17(19.46) |
| rc202 | 214.20 | 186.40 | 859.00(29.34) | 864.76(17.16) |
| rc203 | 235.10 | 188.66 | 926.51(10.38) | 928.59(8.42) |
| rc204 | 266.33 | 400.14 | 1040.73(13.96) | 1039.69(13.62) |
| rc205 | 220.91 | 147.50 | 795.34(17.91) | 790.46(18.13) |
| pr01 | 127.35 | 88.51 | 287.99(3.50)(+) | 284.20(5.64) |
| pr02 | 128.91 | 99.83 | 383.04(3.09) | 380.80(8.81) |
| pr03 | 93.00 | 180.06 | 310.41(5.40) | 311.30(3.33) |
| pr04 | 133.97 | 125.63 | 341.93(8.13) | 341.89(8.40) |
| pr05 | 115.18 | 144.66 | 391.88(15.63) | 391.23(20.75) |
| pr11 | 140.94 | 168.94 | 311.50(12.20) | 305.05(17.73) |
| pr12 | 138.68 | 187.60 | 402.91(6.09) | 405.88(4.66) |
| pr13 | 90.46 | 183.18 | 373.90(11.58) | 369.98(21.23) |
| pr14 | 119.04 | 226.82 | 405.52(18.82) | 404.36(14.74) |
| pr15 | 103.10 | 355.43 | 458.58(10.36) | 461.24(11.11) |

for most instances there is a clear upward trend of the test performance for both algorithms. In addition, both algorithms almost converged after 30 generations, and there was no overfitting. This verifies the effectiveness of the instance rotation during the training process. Also, the convergence curves of the two algorithms are very close to each other, indicating that including the look-ahead features into the terminals did not influence the GP evolution process much.

Then, we conduct some investigation on the importance of features. To this end, we selected some representative instances (r101, c203 and pr03 with $m = 1$), and plot the frequency (total occurrences) of the terminals in the 30 final policies obtained by the 30 runs of BasicGP and LA-GP. Other instances showed similar patterns. Figs. 3 and 4 show

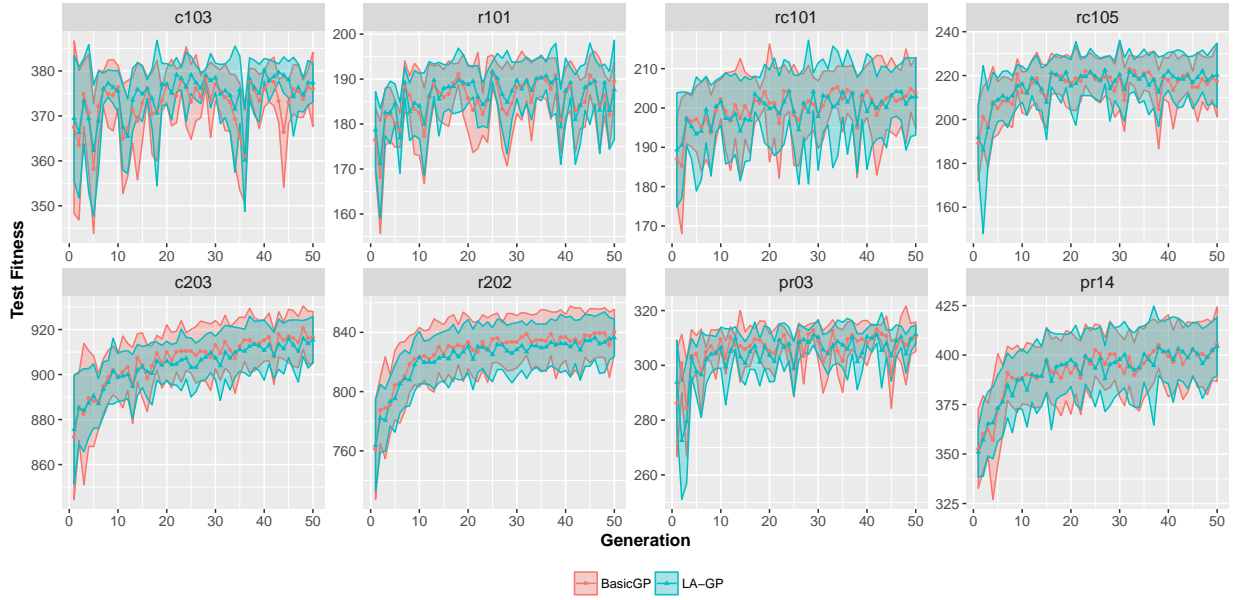


Fig. 2. The convergence curves of the test performance of BasicGP and LA-GP on c103, r101, rc101, rc105, c203, r202, pr03 and pr14.

the terminal frequencies of BasicGP and LA-GP, respectively. Note that in Fig. 3, the frequencies of ANS and MNS are zero, as these two features are not used in BasicGP.

From the figures, one can see that SCORE is always important in all the instances, which is consistent with intuition, as the objective is to maximise the total score. Then, either TSV or TFV or both appeared frequently in the final policies obtained by both algorithms in all the instances. This is also consistent with intuition, as one tends to select POIs with shortest time to start or finish the visit. On the other hand, TO, TC, TA and $SL = TC - TA$ did not often appear. Note that $TSV = \max\{TO, TA\}$, Figs. 3 and 4 indicate that some features are redundant, and TO, TC, TA or/and SL may be removed in the presence of TSV or TFV.

For c203 and pr03, the frequencies of the look-ahead features ANS and MNS are quite low. For r101, ANS and MNS seemed to appear a bit more frequently. More specifically, the frequencies of ANS and MNS in the final policies of the 30 runs of LA-GP is shown in 5. It can be seen that the maximal frequency is ANS appearing for 5 times in runs 9, 20 and 28. Correspondingly, the test fitness of these three runs are 194.41, 190.7 and 190.9, which are higher than the average test fitnesses of both algorithms (189.52 and 187.59). As an example, we investigated the structure of the final policy obtained by run 9. After simplification, the policy can be written as follows.

$$\text{ANS} - \max\{\text{MNS}, \text{BB}\} + 2 \cdot \text{BB} - \text{DUR} - \max\left\{\frac{\max\{\text{DUR}, \text{RemT}\} - \text{TFV}}{\text{TA} - \text{TR}}, \min\{\max\{\text{ANS}, \text{SCORE}\}, \text{RemT} + \text{SL}\}\right\},$$

where $\text{BB} = \text{SCORE} + \text{ANS}$. From the policy, one can see that $\text{SCORE} + \text{ANS}$ can be a promising constructed feature, which

stands for the sum of the score of the current POI and the expected (average) score of next POI. The policy clearly attempts to select the POI with the maximal $\text{SCORE} + \text{ANS}$ value. This shows that the look-ahead features can be helpful to the effectiveness of the policy. Besides, this demonstrates the ability of GP to automatically construct promising features from low-level features.

Finally, Table V shows the average training time (in seconds) of BasicGP and LA-GP over all the 40 TOPTW instances in the experiment. From the table, it is clear that LA-GP is significantly slower than BasicGP. This is mainly due to the calculation of the ANS and MNS features, which requires the enumeration of all the unvisited POIs. In other words, although the look-ahead features provide more comprehensive information for decision making, they are much more time consuming to calculate, which makes the GPHH much less efficient.

V. CONCLUSIONS AND FUTURE WORK

In this paper, a GPHH method was proposed for evolving policies for stochastic TOPTW. In particular, a meta-algorithm for fitness evaluation is developed and a set of features are designed for the terminal set of the GPHH. Two versions of GPHH, one with and the other without the two look-ahead features, are compared on benchmark instances. The results showed that the current terminal set may contain redundant features, and may be improved by identifying and removing the redundant features. Also, it is demonstrated that the look-ahead features can help improve the effectiveness of the policy if used properly. However, simply including the look-ahead features into the terminal set seems not to be effective in employing the look-ahead information. In future work, we will consider feature selection and construction so that the

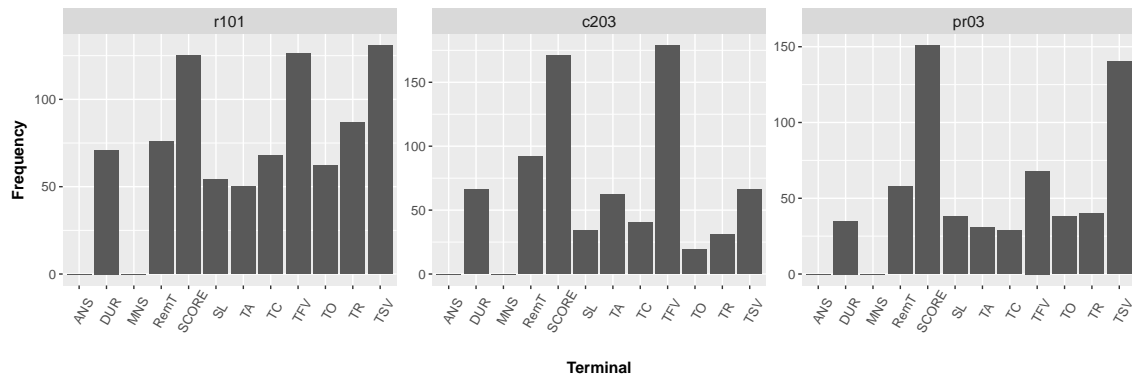


Fig. 3. The terminal frequencies of the final policies obtained by the 30 runs of BasicGP on r101, c203 and pr03 with $m = 1$.

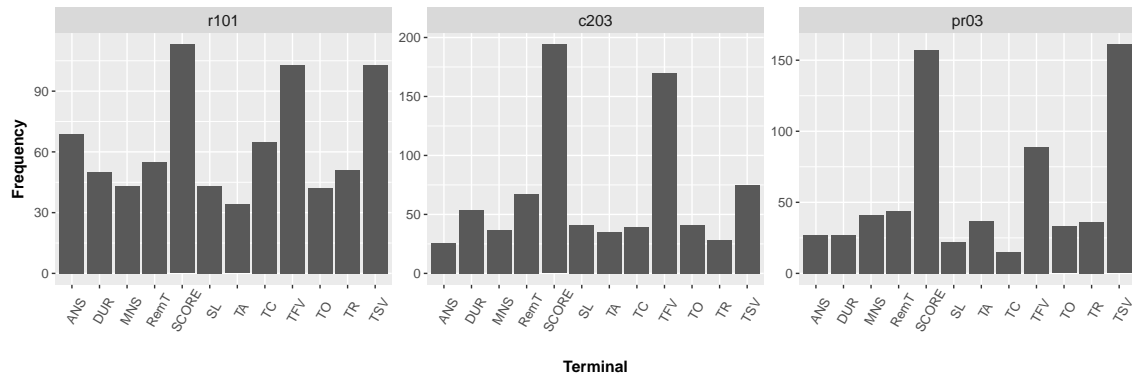


Fig. 4. The terminal frequencies of the final policies obtained by the 30 runs of LA-GP on r101, c203 and pr03 with $m = 1$.

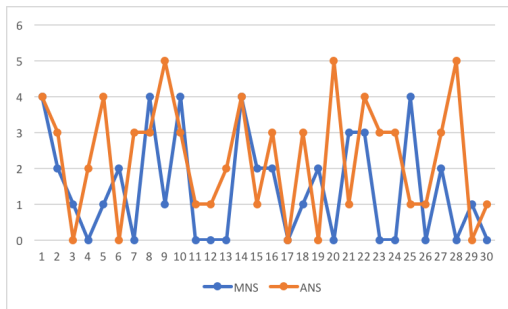


Fig. 5. The frequencies of ANS and MNS in the final policies obtained by the 30 runs of LA-GP on r101 with $m = 1$.

look-ahead information can be employed in a more intelligent way. In addition, new features with better balance between effectiveness and efficiency are to be developed.

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TABLE IV

THE TEST FITNESS OF THE COMPARED GPHH METHODS ON THE INSTANCES WITH $m = 3$. IF BASICGP IS STATISTICALLY SIGNIFICANTLY BETTER THAN LA-GP UNDER WILCOXON'S RANK SUM TEST WITH $\alpha = 0.05$, THEN THE CORRESPONDING RESULT IS MARKED WITH "(+)". IF BASICGP IS STATISTICALLY SIGNIFICANTLY WORSE THAN LA-GP, THEN THE CORRESPONDING RESULT IS MARKED WITH "(-)".

| Instance | SOS | SOT | BasicGP | LA-GP |
|----------|---------|---------|------------------|----------------|
| c101 | 488.46 | 339.02 | 759.64(9.63) | 764.59(4.83) |
| c102 | 538.34 | 386.50 | 854.93(8.62) | 858.90(11.82) |
| c103 | 576.90 | 503.08 | 908.82(6.69) | 909.86(1.38) |
| c104 | 668.90 | 612.58 | 954.18(13.46) | 953.51(7.83) |
| c105 | 527.74 | 345.38 | 805.41(9.47) | 806.09(8.17) |
| r101 | 240.84 | 181.61 | 434.46(11.84) | 439.18(8.95) |
| r102 | 297.38 | 285.34 | 603.07(15.47) | 602.76(21.14) |
| r103 | 323.44 | 270.31 | 656.17(11.44) | 653.98(16.79) |
| r104 | 306.39 | 447.93 | 711.40(15.70) | 713.47(10.82) |
| r105 | 256.52 | 202.53 | 546.63(16.83) | 548.51(13.70) |
| rc101 | 275.79 | 187.65 | 562.82(12.49) | 563.28(16.79) |
| rc102 | 269.82 | 243.99 | 620.66(13.54) | 618.41(17.62) |
| rc103 | 279.12 | 205.06 | 667.10(20.07) | 665.63(15.06) |
| rc104 | 312.69 | 557.12 | 721.00(14.32) | 721.39(11.40) |
| rc105 | 250.46 | 246.25 | 611.26(13.20)(-) | 618.42(9.23) |
| c201 | 542.52 | 722.00 | 1765.92(5.04)(+) | 1757.61(4.48) |
| c202 | 675.46 | 616.58 | 1774.19(5.69) | 1773.14(10.79) |
| c203 | 898.84 | 622.68 | 1767.32(5.81) | 1767.91(5.43) |
| c204 | 1008.46 | 948.68 | 1779.52(7.46) | 1775.30(16.58) |
| c205 | 659.62 | 599.82 | 1782.96(3.33)(-) | 1784.29(2.61) |
| r201 | 588.73 | 402.01 | 1357.85(20.99) | 1350.16(20.94) |
| r202 | 786.46 | 593.48 | 1415.45(8.58)(+) | 1408.21(8.22) |
| r203 | 750.60 | 858.52 | 1438.33(1.79) | 1437.19(3.14) |
| r204 | 1028.85 | 1008.51 | 1440.19(3.61) | 1438.27(13.71) |
| r205 | 788.44 | 819.42 | 1437.14(4.51) | 1437.14(3.19) |
| rc201 | 640.44 | 407.00 | 1597.26(17.15) | 1595.38(15.60) |
| rc202 | 768.40 | 524.57 | 1656.74(20.01) | 1653.82(22.96) |
| rc203 | 703.26 | 599.21 | 1712.06(4.90)(+) | 1708.82(7.67) |
| rc204 | 793.36 | 1299.10 | 1720.24(2.49) | 1720.54(1.17) |
| rc205 | 622.49 | 555.27 | 1624.24(22.86) | 1629.21(31.23) |
| pr01 | 246.86 | 294.71 | 559.84(12.25)(-) | 568.47(6.11) |
| pr02 | 338.35 | 474.12 | 864.08(10.97) | 853.41(21.48) |
| pr03 | 266.48 | 383.93 | 770.33(19.91) | 775.33(8.95) |
| pr04 | 391.48 | 448.86 | 855.00(4.76) | 846.50(41.50) |
| pr05 | 276.33 | 464.83 | 908.18(21.45) | 911.95(14.53) |
| pr11 | 369.38 | 457.60 | 621.10(8.69) | 619.72(14.64) |
| pr12 | 396.00 | 657.75 | 914.26(10.34) | 912.31(13.59) |
| pr13 | 341.32 | 611.21 | 889.10(16.95)(+) | 878.93(22.28) |
| pr14 | 385.58 | 656.05 | 920.71(21.19) | 929.29(16.20) |
| pr15 | 401.63 | 739.16 | 970.11(23.57) | 974.61(18.28) |

TABLE V

THE AVERAGE TRAINING TIME (IN SECONDS) OF BASICGP AND LA-GP OVER ALL THE TOPTW INSTANCES.

| | BasicGP | LA-GP |
|-----------|---------|----------|
| Avg. Time | 966.30 | 26182.05 |

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